

The Road from Leibniz to Turing, Part I

DA2205, DD3001

October 4, 2012

Leibniz's Dream

Source: Lecture material is based on *The Universal Computer* by Martin Davis

- **Born:** Leipzig, Germany, July 1646.
- **Died:** Hanover, Germany, Nov 1716.
- **Father:** Professor of Moral Philosophy, University of Leipzig.

Gottfried had unrestricted access to his father's extensive library.



- **Education:**

- ★ BA in Philosophy, University of Leipzig, 1662
- ★ Master's in Philosophy, Univ. of Leipzig, 1664
- ★ Law degree, University of Leipzig, 1665
- ★ Doctorate in Law, University of Altdorf, 1666



- **Employment:**

Wealthy noble patrons

- ★ Baron von Boyneburg, 1666 – 1674
 - Diplomatic missions for Elector of Mainz
 - Got to spend time in Paris.
- ★ Dukes of Hanover, 1675 – 1716
 - Political adviser, historian, librarian



Gottfried Leibniz: Major Research Achievements

- Prominent figure in the history of mathematics and the history of philosophy.
- Infinitesimal calculus – probably independently of Newton!
We still use his notation today.
- Towards the development of computation
 - ★ Invented mechanical calculator capable of
 $+$, $-$, \times , \div
 - ★ Contributions to Formal Logic – unpublished in lifetime
 - ★ Believed human reasoning could be reduced to calculations
 - ★ Envisaged a *calculus ratiocinator* – resembling symbolic logic – to make such calculations feasible.
 - ★ Studied binary notation



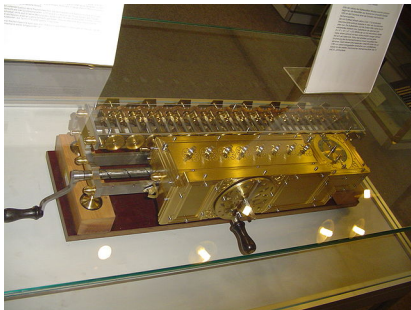
Gottfried Leibniz: Fame & fortune during lifetime

- Much.
- Was a major courtier to a powerful German royal dynasty → good lifestyle but forced to perform time sapping duties unrelated to his interests – genealogy.
- Corresponded with the major thinkers of his day.
- Made a member of the Royal Society of London in 1673.
- However, reputation towards end of life was in decline.
- Especially tarnished by the controversy with Newton over the discovery of calculus.
- Posthumously though his reputation was restored.

- As a teenager was introduced to the work of Aristotle.
- This inspired a "*wonderful idea*":
 - Seek an alphabet whose elements represent concepts
 - This alphabet would form a language
 - In this language by symbolic reasoning determine
 - which sentences in the language were true **and**
 - what logical relationships existed among them.
- Leibniz held onto this vision throughout his lifetime...
- ... and made some progress towards it.

- From 1672-1676 Leibniz was in Paris on a diplomatic mission.
- During this time was exposed to the modern Mathematics of the day which had been fueled by the
 - ★ systemization of the techniques for dealing with algebraic expressions
 - ★ realization geometry could be expressed as algebra.
- Made contact with the great thinkers of the time.
- Own research:
 - ★ Leibniz series for π
 - ★ many of the concepts and ideas needed for his derivation of calculus.
- Convinced himself that it is crucial to have an appropriate symbolism when representing and solving problems.

The Stepped Reckoner



- **1671:** Leibniz began work on the "*Stepped Reckoner*" a machine that could $+$, $-$, \times , \div .
- Prototypes made in Hanover by a craftsman working under Leibniz's supervision.
- Not an unambiguous success as did not fully mechanize the operation of carrying. But its "Leibniz wheel" was a success.

Leibniz saw three strands to his problem:

- Create a compendium of all human knowledge - Crazy then!
Crazy now?
- Identify key underlying notions in this compendium and provide them with appropriate symbols
- Rules of deduction encoded as manipulation of these symbols
 - Leibniz's **calculus ratiocinator**, the algebra of logic.

Leibniz's attempt at symbolic logic

DEFINITION 3. *A is in L, or L contains A, is the same as to say that L can be made to coincide with a plurality of terms taken together of which A is one. $B \oplus N = L$ signifies that B is in L and that B and N together compose or constitute L. The same thing holds for a larger number of terms.*

AXIOM 1. $B \oplus N = N \oplus B$.

POSTULATE. Any plurality of terms, as A and B, can be added to compose a single term $A \oplus B$.

AXIOM 2. $A \oplus A = A$.

PROPOSITION 5. *If A is in B and $A = C$, then C is in B.*

For in the proposition *A is in B* the substitution of A for B gives *C is in B*.

PROPOSITION 6. *If C is in B and $A = B$, then C is in A.*

For in the proposition *C is in B* the substitution of A for B gives *C is in A*.

PROPOSITION 7. *A is in A.*

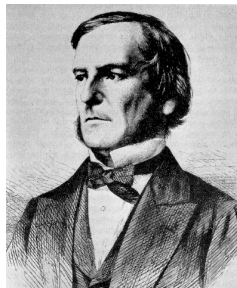
For *A is in $A \oplus A$* (by Definition 3). Therefore (by Proposition 6) *A is in A*.

.....
PROPOSITION 20 *If A is in M and B is in N, then $A \oplus B$ is in $M \oplus N$.*

**Boole turns logic
into algebra**

- **Born:** Lincoln, England, Nov 1815.
- **Died:** Cork, Ireland, Dec 1864.
- **Father:** Cobbler and an inept businessman

From the age of 16 George was responsible for providing financially for the family.



George Boole: Education and Employment

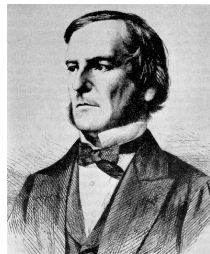
- **Education:**

- ★ Elementary school education
- ★ Self-taught with some guidance from the *Lincoln Mechanics' Institution*

- **Employment:**

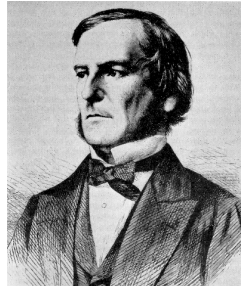
- ★ School teacher, ~ Lincoln, 1832–1835
- ★ Ran and taught in schools he founded, Lincoln, 1835–1849
- ★ Professor, University College Cork, Ireland, 1849–1864

Note: Before his professorship he became an active mathematician while running his school !



George Boole: Major Research Achievements

- Boolean Logic - the basis of calculations in the modern digital computer !
- In the *The Laws of Thought* demonstrated that logical deduction could be seen as a branch of mathematics (algebra).
- Also made contributions to differential equations.



- Some.
- He received a medal from the *Royal Society* for a publication on linear differential equations.
- Made it to Professor of Mathematics though in a provincial backwater !
- May have progressed further and earlier if he had had a more conventional and privileged background.

Boole was aware that the power of algebra is derived from

- the fact that it has the symbols representing both quantities and operations **and**
- these obey a small number of basic rules or laws.

Classical logic and introduction of symbols

- Classical logic involves sentences such as
 1. All plants are alive.
 2. No hippopotamus is intelligent.
 3. Some people speak English.
- Boole realized that if we reason about the words alive, hippopotamus, or people what is significant for each is the class of all individuals described by the word.
- Boole saw how this reasoning could be expressed in terms of an algebra of such classes.
- Boole represented classes by letters. In his own words...

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Introduction of symbols for classes

*“... If an adjective, as '**good**', is employed as a term of description, let us represent by a letter, as y , all things to which the description '**good**' is applicable, i.e. '**all good things**', or the class '**good things**'. Let it further be agreed, that by the combination xy shall be represented that class of things to which the names of descriptions represented by x and y are simultaneously applicable. Thus, if x alone stands for '**white things**' and y for '**sheep**', let xy stand for '**white sheep**' and in like manner, if z stands for '**horned things**', ...let zxy represent '**horned white sheep**'.*

– George Boole

Introduction of symbols for classes

- Following Boole's example and notation, what is yy ?
- yy is the class of sheep that are also ...sheep. Therefore

$$yy = y$$

- Most of Boole's system of logic is based on the fact
when x stands for a class, the equation $xx = x$ is always true

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when x stands for a class, the equation $xx = x$ is always true

- In ordinary algebra, where x stands for a number, when is the equation $xx = x$ true?
- **Answer:** The equation is true when either $x = 0$ or $x = 1$.
- Boole concluded
algebra of logic \equiv ordinary algebra restricted to the numbers $\{0, 1\}$

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To have an **algebra for logic** similar to **ordinary logic**, need

1. Definition for classes of the binary operators

★ Addition

$x + y \equiv$ class consisting of both the elements of x and y

★ Subtraction

$x - y \equiv$ class consisting of the elements in x not in y

2. Identity elements for each binary operator

★ **Multiplication**: Need 1 s.t. $1x = x$. 1 is then the class

"containing every object under consideration" $\equiv 1$

★ **Addition**: Need a 0 s.t. $x + 0 = x$. 0 is then the class

"containing nothing" $\equiv 0$

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"containing nothing" $\equiv 0$

- In this logic what is

- ★ $0x$?

- ★ $1 - x$?

- ★ $x(1 - x)$?

- Note the final expression can be derived from

$$xx = x \implies x - xx = 0 \implies x(1 - x) = 0$$

- In this logic
 - ★ $0x = 0$
 - ★ $1 - x =$ class "*containing every object not in x* "
 - ★ $x(1 - x) = 0$ i.e. nothing can belong **and** not belong to a class
- Note the final expression can be derived from

$$xx = x \implies x - xx = 0 \implies x(1 - x) = 0$$

- Aristotle's logic focused on inferences of a special type called **syllogisms**.
- Inference is from a pair of **premises** to a **conclusion**.
- The premises and conclusions must be representable by a sentence of this form

Sentence type

All X are Y No X are Y Some X are Y Some X are not Y

- A **valid** syllogism is

$$(\text{All } X \text{ are } Y) \text{ and } (\text{All } Y \text{ are } Z) \implies (\text{All } X \text{ are } Z)$$

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- Translation of premises into Boole's algebra

$$X = XY \quad \text{and} \quad Y = YZ$$

then the conclusion is

$$X = XY = X(YZ) = (XY)Z = XZ$$

- Translation of conclusion back to words is (All X are Z)

- Of course not all syllogism are valid.

- An **invalid** syllogism is

$$(\text{All } X \text{ are } Y) \text{ and } (\text{All } X \text{ are } Z) \implies (\text{All } Y \text{ are } Z)$$

- This invalid syllogism cannot be derived in Boole's logic.
- Therefore Boole's logic includes Aristotle's logic, but is capable of reasoning far beyond it.

- ✓ It is easy to use Boole's algebra as a system of rules for calculating \implies provided the calculus ratiocinator.
- ✓ Boole showed logical deduction could be developed as a branch of mathematics.
- ✓ Boole's system of logic included Aristotle's logic and went beyond it.

- ✗ Boole's logic was still fairly crude. Cannot represent statements of the type:

All failing students are either stupid or lazy

- ✗ Boole's logic did not encompass his logic as a fully-fledged deductive system in which all the rules are deduced from a small set of axioms.

Frege: From breakthrough to despair

- **Born:** Wismar, Germany, Nov 1848.
- **Died:** Bad Kleinen (near Wismar), Germany July 1925.
- **Father:** Co-founded and was headmaster of a girls' high school
- **Mother:** Ran the school after Frege's father's death.



- **Education:**

- ★ University of Jena, 1869 – 1870
- ★ University of Göttingen, 1871 – 1873
- ★ Awarded PhD in Mathematics in 1873



- **Employment:**

- ★ Unpaid lecturer, University of Jena, 1874 – 1879
- ★ Associate Professor, University of Jena, 1879 – 1918



- Not much !
- His achievements were mainly unrecognized in his lifetime.
- Didn't make it to Professor.
- Left broken by his work and embittered at the time of his death.

Gottlob Frege: Major Research Achievements

- Modern logic - axiomatic predicate logic !
- Introduced and discovered how to manipulate the quantifiers (\forall, \exists), truth functions (\neg, \wedge, \vee and \implies), variables and predicates.
- Developed artificial language with precise rules of grammar \implies logical inferences as purely mechanical operations.
- This logic system encompassed all of the reasoning used by mathematics.
- Frege's logic was an enormous advance over Boole's



Introducing his logical system



In 1879 Frege published

Begriffsschrift - Concept Notation, the Formal Language of the Pure Thought like that of Arithmetics.

outlining his logic system.

Statements in Frege's Logic – Universal quantifier

The sentence

All horses are mammals.

in can be expressed in “Frege speak”

If x is a horse, **then** x is a mammal.

↓

$(\forall x)(\mathbf{if} \ x \text{ is a horse, then } x \text{ is a mammal}).$

↓

$(\forall x)(x \text{ is a horse} \supset x \text{ is a mammal}).$

↓

$(\forall x)(\text{horse}(x) \supset \text{mammal}(x)).$

- These steps of abstract are possible because
 - ★ The original statement is true for all x – symbol \forall denotes “for all”.
 - ★ Logical relation **if ... , then ...** is symbolized by \supset .

Statements in Frege's Logic – Existential quantifier

The sentence

Some horses are pure-bred

in can be expressed in “Frege speak”

x is a horse **and** x is pure-breed.

↓

$(\exists x)(x \text{ is a horse and } x \text{ is pure-bred}).$

↓

$(\exists x)(x \text{ is a horse} \wedge x \text{ is pure-bred}).$

↓

$(\exists x)(\text{horse}(x) \wedge \text{pure-bred}(x)).$

- These steps of abstract are possible because
 - ★ Original statement is only true for **some** x - \exists denotes “there exists”.
 - ★ Relation **and** is symbolized by \wedge .

Symbol	Interpretation
\neg	not ...
\vee	... or ...
\wedge	... and ...
\supset	if ..., then ...
\forall	every
\exists	some

- Boole's logic could not express the statement
All failing students are either stupid or lazy
- How do we express this assertion in Frege's system?

$$(\forall x)(\text{Failing}(x) \supset (\text{Stupid}(x) \vee \text{Lazy}(x)))$$

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Most fundamental rule of inference is as follows

- \triangle is a sentence
- \diamond is a sentence
- If both

\triangle **and** $(\triangle \supset \diamond)$

are true then \diamond is true.

- **Question:**

Why did Frege develop this logic system?

- **Answer:** Frege believed mathematics is nothing but logic.

Wanted to show

- ★ given his logic and the concept of set, then all of mathematics follows
- ★ every concept in mathematics can be explicitly defined in terms of logic **and**
- ★ every statement in mathematics can be translated by a well-formed formula of logic.
- ★ all of the basic principles of mathematics could be derived from the fundamental laws of logic.

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- ★ all of the basic principles of mathematics could be derived from the fundamental laws of logic.

- Given the natural numbers can derive much of mathematics

“God made natural numbers, all the rest are made by Man.

– Leopold Kroneker (1823 - 1891)

- Therefore, Frege was going to achieve his goal by providing a purely logical theory of the natural numbers - $\{1, 2, 3, \dots\}$
- He outlined his project in the book *The Foundations of Arithmetic*, 1884.
- He developed his project in detail in the two-volume set *The Fundamental Laws of Arithmetic*. Vol I (1893), Vol II (1903).

- To fulfil his goal Frege introduced the concept of a set.
- If b is an element of a set a then write

$$b \in a$$

- Added two principles to basic logic:
 - ★ Sets with the same members are the same set:

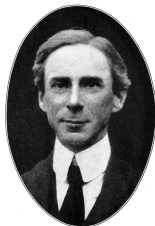
$$x = y \iff \forall z(z \in x \iff z \in y)$$

- ★ Given any property F there is a set consisting of all those things that had F

$$(\exists y)(\forall x)(x \in y \iff F(x))$$

Disaster - A letter in from Bertrand Russell

- In 1902 as his book was going to press Frege received a letter from the young British philosopher Bertrand Russell.



- Russell brought Frege's attention to the axiom of set existence

$$(\exists y)(\forall x)(x \in y \iff F(x))$$

and asked what happens when $F(x)$ represents the property

"x is not a member of itself"

, that is $F(x) \iff \neg(x \in x)$ and x is a set. **What happens?**

- According to Frege there should be a set consisting of all and only those sets that don't belong to themselves:

$$(\exists y)(\forall x)(x \in y \iff \neg(x \in x))$$

- But if this is true for all x that

$$x \in y \iff \neg(x \in x)$$

then it's true in particular for y :

$$y \in y \iff \neg(y \in y)$$

A contradiction !

- Thus Frege's **Basic principle of logic** was not true.
- A decade of his work was invalidated.

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“There is nothing worse can happen to a scientist than to have the foundation collapse just as the work is finished. I have been placed in this position by a letter from Mr. Bertrand Russell.”

– Gottlob Frege, Appendix of The Fundamental Laws of Arithmetic, 2nd volume

- Frege never really recovered from this blow to his work.

- ✓ *Begriffsschrift* can be seen as embodying the universal language of logic envisioned by Leibniz.
- ✓ *Begriffsschrift* encapsulated the logic used in ordinary maths \implies mathematical activity could be investigated by mathematical methods.
- ✗ Frege's logic not efficient for calculations.
- ✗ Frege's logic provides no procedures for determining whether some desired conclusion can be deduced from a set of premises.

Cantor: Detour through infinity

- Cantor's work marks
the beginning of the death of certainty in mathematics and
the birth of computer science !
- Statement based on the fact

Cantor's work raised troubling paradoxes



Hilbert & Gödel worked to resolved them



Turing was inspired in turn by this work



John von Neumann borrowed from Turing to design the EDVAC.

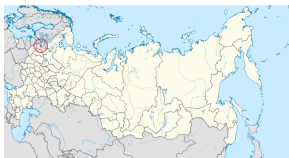
He realized the computing machine is just a logic machine.

- **Born:** St. Petersburg, Russia, 1845.

Family moved to Germany when Georg was 11.

- **Died:** Halle, Germany, January 1918.
- **Father:** Successful businessman.

Inherited sufficient money on father's death to pursue an academic career.



- **Education:**

- ★ University of Zürich, 1862
- ★ University of Berlin, 1863–1867
- ★ University of Göttingen, Summer 1866
- ★ Awarded PhD in Mathematics in 1867

- **Employment:**

- ★ Teacher in a girl's school, Berlin, 1868
- ★ University of Halle, Germany
 - Privatdozent, 1869–1872
 - Extraordinary Professor, 1872–1879
 - Professor, 1879–1913



Georg Cantor: Major Research Achievements

- Inventor of set theory
- Established the importance of one-to-one correspondence between the members of two sets
- Created a profound and coherent mathematical theory of the infinite.

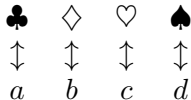


Georg Cantor: Fame & fortune during lifetime

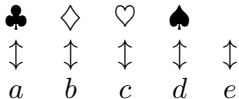
- Some, but had to endure much criticism !
- Full Professor by the age of 34.
- Work not accepted by much of the establishment:
 - ★ “grave disease” infecting mathematics, Poincaré
 - ★ “scientific charlatan”, “corrupter of youth”, Kronecker
 - ★ “utter nonsense”, “laughable”, Wittgenstein
- Thus blocked from Professorships at prestigious universities.
- Did have supporters: Dedekind, Weierstrass and Mittag-Leffler.
Work was “... about one hundred years too soon.”, Mittag-Leffler
- Received prestigious accolades later on though.
 - ★ 1904 - *Sylvester Medal* from the Royal Society
 - ★ 1911 - Honorary Doctorate from St. Andrews University, Scotland

Deciding if two sets have the same size

- Two sets have the same number of members (cardinality) if the members in each set can be matched up in a 1-1 fashion.
- Sets $\{\clubsuit, \diamondsuit, \heartsuit, \spadesuit\}$ and $\{a, b, c, d\}$ have same cardinality as:



- Sets $\{\clubsuit, \diamondsuit, \heartsuit, \spadesuit\}$ and $\{a, b, c, d, e\}$ do not as:



- Cantor applied the idea of 1-1 matching with infinite sets.

Infinite sets of the same cardinality

- Consider these two sets

- ★ set of all natural numbers, $1, 2, 3, 4, \dots$ and

- ★ set of all even natural numbers $2, 4, 6, \dots$

Do these two sets have the same cardinality?

- **Yes**

1	2	3	4	...
↕	↕	↕	↕	
2	4	6	8	...

- Cantor investigated which other infinite sets could be matched in 1-1 correspondence...

Infinite sets of the same cardinality

- Is the set of rational numbers larger than the set of natural numbers ?

- No

- Can list all the possible fractions as follows

$$\left| \frac{1}{1} \right| \left| \frac{1\ 2}{2\ 1} \right| \left| \frac{1\ 2\ 3}{3\ 2\ 1} \right| \left| \frac{1\ 2\ 3\ 4}{4\ 3\ 2\ 1} \right| \left| \frac{1\ 2\ 3\ 4\ 5}{5\ 4\ 3\ 2\ 1} \right| \dots$$

Grouped so numerator + denominator equals 2, 3, 4, 5, ...

- Matching up is now trivial

$$\begin{array}{ccccccc} \frac{1}{1} & \frac{1}{2} & \frac{2}{1} & \frac{1}{3} & \frac{2}{2} & \frac{3}{1} & \frac{1}{4} & \dots \\ \updownarrow & \updownarrow & \updownarrow & \updownarrow & \updownarrow & \updownarrow & \updownarrow & \\ 1 & 2 & 3 & 4 & 5 & 6 & 7 & \dots \end{array}$$

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$$\left| \frac{1}{1} \right| \left| \frac{1}{2} \frac{2}{1} \right| \left| \frac{1}{3} \frac{2}{2} \frac{3}{1} \right| \left| \frac{1}{4} \frac{2}{3} \frac{3}{2} \frac{4}{1} \right| \left| \frac{1}{5} \frac{2}{4} \frac{3}{3} \frac{4}{2} \frac{5}{1} \right| \dots$$

Grouped so numerator + denominator equals 2, 3, 4, 5, ...

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Infinite sets can come in different sizes

- Is the set of real numbers larger than the set of natural numbers ?
- **Yes**
- Note though the set of algebraic numbers has the same size as the natural numbers.
- Cantor let
 - ★ \aleph_0 represent the cardinality of the set of natural numbers
 - ★ \mathcal{C} the cardinality of the set of real numbers.

Ranking of members of a set

- Members of $\{\clubsuit, \diamond, \heartsuit\}$ can be ranked in 6 different ways:

1 st	2 nd	3 rd	1 st	2 nd	3 rd	1 st	2 nd	3 rd	1 st	2 nd	3 rd	1 st	2 nd	3 rd	1 st	2 nd	3 rd
↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓
♣	◇	♥	♣	♥	◇	◇	♣	♥	◇	♥	♣	♥	♣	◇	♥	◇	♣

For each ranking, easy to label the rank of a set member.

- But what about infinite sets? Consider the natural numbers.
- Can obviously list these in any order we like.
- Say all even numbers are listed and then all odd numbers

2, 4, 6, ..., 1, 3, 5, ...

What is the rank of number '2' ?

Ranking of members of a set

- Members of $\{\clubsuit, \diamond, \heartsuit\}$ can be ranked in 6 different ways:

1 st	2 nd	3 rd	1 st	2 nd	3 rd	1 st	2 nd	3 rd	1 st	2 nd	3 rd	1 st	2 nd	3 rd	1 st	2 nd	3 rd
↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓	↓
♣	◇	♥	♣	♥	◇	◇	♣	♥	◇	♥	♣	♥	♣	◇	♥	◇	♣

For each ranking, easy to label the rank of a set member.

- But what about infinite sets? Consider the natural numbers.
- Can obviously list these in any order we like.
- Say all even numbers are listed and then all odd numbers

$2, 4, 6, \dots, 1, 3, 5, \dots$

What is the rank of number '1' ?

Ranking of members of an infinite set

- Cantor introduced the first *transfinite ordinal number* ω .

1 st	2 nd	3 rd	...	$(\omega + 1)^{\text{th}}$	$(\omega + 2)^{\text{th}}$	$(\omega + 3)^{\text{th}}$...
↓	↓	↓		↓	↓	↓	
2	4	6	...	1	3	5	...

- Natural numbers can be ranked in many different ways using larger and larger transfinite ordinal numbers.
- Set {all ordinal numbers needed to define all rankings of \mathbb{N} } said to have cardinality \aleph_1 .
- Cantor proved $\aleph_1 > \aleph_0$ and there is no cardinal number \aleph s.t. $\aleph_0 < \aleph < \aleph_1$.
- But it doesn't stop here....

An infinite number of infinite cardinalities

- What about ranking the members of sets of cardinality \aleph_1 ?
- For these rankings need to introduce the transfinite ordinal number ω_1
- Set {all ordinal numbers needed to define all rankings of sets with cardinality \aleph_1 } has cardinality \aleph_2 .
- There is no end to this process. Can define $\aleph_3, \aleph_4, \dots, \aleph_\omega, \dots$

- Cantor used versions of the diagonal method to show
 - ★ The cardinality of the set of real numbers is larger than \aleph_0
 - ★ the cardinality of a set S is less than the cardinality of the power set of S .
- The players in this story also used the diagonal method
 - ★ Russell when considering the *set of all sets*.
 - ★ Gödel proving his first incompleteness theorem
 - ★ Turing in analyzing the *Entscheidungsproblem*.

The diagonal method overview

E_0	=	m	m	m	m	m	m	m	m	m	m	m	...
E_1	=	w	w	w	w	w	w	w	w	w	w	w	...
E_2	=	m	w	m	w	m	w	m	w	m	w	m	...
E_3	=	w	m	w	m	w	m	w	m	w	m	w	...
E_4	=	w	m	m	w	m	m	w	m	w	m	w	...
E_5	=	m	w	m	w	w	m	w	m	w	m	w	...
E_6	=	m	w	m	w	w	m	w	m	w	m	w	...
E_7	=	w	m	m	w	m	w	m	w	m	w	m	...
E_8	=	m	m	w	m	w	m	w	m	w	m	w	...
E_9	=	w	m	w	m	m	w	w	m	w	m	w	...
E_{10}	=	w	w	m	w	m	w	m	w	m	w	m	...
E_{11}	=	m	w	m	w	w	m	w	m	m	w	m	...
\vdots		\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\vdots	\ddots
E_u	\neq	w	m	w	w	m	w	m	m	m	m	w	

- Cantor's diagonal method for the existence of uncountable sets.
- Bottom sequence cannot occur anywhere in the list of sequences above.
- Therefore cannot list all the infinite sequences of the above form.

- Valid reasoning with Cantor's transfinite can lead to paradoxes.

- ★ **1895**: Cantor considered

- What is the cardinality of the set of all cardinal numbers ?

- ★ **1897**: Burali-Forti published a similar paradox when considering the set of all transfinite ordinal numbers.

- Bertrand Russell while considering Cantor's work asked

- Can there be a set of all sets?*

- led to paradox of sets who are members of themselves

- letter to Frege.

- Mathematics in crisis.

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